קודם: שקפים 1-4 על הלוח (הדוגמא גם על שקף 4-5)
פירוק LU
$\left.\begin{array}{c}\text { A } \\ \left(\begin{array}{cc}0.0003 & 3 \\ 1 & 1\end{array}\right)\end{array} \begin{array}{c}\mathbf{X} \\ x_{1} \\ x_{2}\end{array}\right)=\left(\begin{array}{c}\mathbf{B} \\ 2.0001 \\ 1\end{array}\right) \begin{gathered}\text { (משיעור שעבר): }\end{gathered}$ הפירוק (עם החלפת השורות):
L U
A'
$\left(\begin{array}{cc}1 & 0 \\ .0003 & 1\end{array}\right) \cdot\left(\begin{array}{cc}1 & 1 \\ 0 & 2.9997\end{array}\right)=\left(\begin{array}{cc}1 & 1 \\ 0.0003 & 3\end{array}\right)$
$\left(\begin{array}{cc}\mathbf{L} \\ 1 & 0 \\ .0003 & 1\end{array}\right) \cdot\left(\begin{array}{c}\mathbf{Y} \\ y_{1} \\ y_{2}\end{array}\right)=\left(\begin{array}{c}\mathbf{B} \\ 1 \\ 2.0001\end{array}\right)$
שלב 1:

שלב 2:
$\left(\begin{array}{cc}1 & 1 \\ 0 & 2.9997\end{array}\right) \cdot\binom{x_{1}}{x_{2}}=\binom{1}{1.9998}$
void ludcmp(float **a, int n , int *indx, float *d)
Given a matrix a[1..n][1..n], this routine replaces it by the $L U$ decomposition of a rowwise permutation of itself. a and $\mathbf{n}$ are input. a is output, arranged as in equation (2.3.14) above; indx[1..n] is an output vector that records the row permutation effected by the partial pivoting; $d$ is output as $\pm 1$ depending on whether the number of row interchanges was even or odd, respectively. This routine is used in combination with lubksb to solve linear equations or invert a matrix.
$a[1 \ldots n][1 \ldots n] \Rightarrow\left(\begin{array}{lll}\beta_{11} & \beta_{12} & \beta_{13} \\ \alpha_{21} & \beta_{22} & \beta_{23} \\ \alpha_{31} & \alpha_{32} & \beta_{33}\end{array}\right) \begin{gathered}\text { מחריף הרצה: a }\end{gathered}$
indx[1...n] שומר את החלפות השורות:
$d= \pm 1$ שומר אם היה מoפר זוגי (1+) או אי-דוגי (1-) של החלפות שורות:
void lubksb(float **a, int n, int *indx, float b[])
Solves the set of $n$ linear equations $A \cdot X=B$. Here $a[1 . . n][1 . . n]$ is input, not as the matrix $A$ but rather as its $L U$ decomposition, determined by the routine ludcmp. indx[1..n] is input as the permutation vector returned by ludcmp. $b[1 . . n]$ is input as the right-hand side vector $B$, and returns with the solution vector $X$. $a$, $n$, and indx are not modified by this routine and can be left in place for successive calls with different right-hand sides $b$. This routine takes into account the possibility that $b$ will begin with many zero elements, so it is efficient for use in matrix inversion.

$$
\text { b[1. . .n] : B } \rightarrow \mathrm{X} \text { אחרי הרצה: b } \mathrm{b} \text { אתריף }
$$

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float **a,*b,d;
int n,*indx;
ludcmp(a,n,indx,&d);
lubksb(a,n,indx,b);
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## SVD פירוק

void svdcmp(float **a, int m, int $n$, float w[], float **v)
Given a matrix a[1..m][1..n], this routine computes its singular value decomposition, $A=U \cdot W \cdot V^{\top}$. The matrix $U$ replaces a on output. The diagonal matrix of singular values $W$ is output as a vector w[1..n]. The matrix $V$ (not the transpose $V^{T}$ ) is output as $\mathrm{v}[1 . . \mathrm{n}][1 . . \mathrm{n}]$.
$a[1 . . . m][1 \ldots n]: A \Rightarrow U \quad$ מחליף את a אחרי הרצה:
$w[1 \ldots n]: \mathrm{W}=\left(\begin{array}{lll}w_{1} & & 0 \\ { }^{2}{ }_{2} & \\ 0 & \ddots & w_{n}\end{array}\right) \quad v[1 \ldots n][1 \ldots n]: V$
אז, מאפסים את ה- w הקטנים, ואחרי זה קוראים (אפשר עם Oדרה של b ) ל:
void subksb(float ${ }^{* *} u$, float $w[]$, float ${ }^{* *} v$, int $m$, int $n$, float b[], float x[]$)$
Solves $A \cdot X=B$ for a vector $X$, where $A$ is specified by the arrays $u[1 . . \mathrm{m}][1 . . \mathrm{n}]$, $\mathrm{w}[1 . . \mathrm{n}], \mathrm{v}[1 . \mathrm{n}][1 . . \mathrm{n}]$ as returned by svdemp. m and n are the dimensions of a , and will be equal for square matrices. $\mathrm{b}[1 . . \mathrm{m}]$ is the input right-hand side. $\mathrm{x}[1 . \mathrm{n}]$ is the output solution vector. No input quantities are destroyed, so the routine may be called sequentially with different b's.

